

Modified Algorithm for Shortening Path Kruskal Algorithm to Communicate Between the Patient and The Ambulance During COVID-19

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Abstract. The COVID-19 virus has highlighted the critical need for efficient and rapid communication between patients in distress and emergency medical services, especially ambulance dispatch systems. The following research paper presents a adapted algorithm intended to adjust the Kruskal method for shortening pathways in communication networks linking patients and ambulances through critical situations, confirming timely reply and effective resource distribution. The suggested method takes into reflection the unique challenges posed by COVID-19, as well as the essential to minimize patient-ambulance interaction to decrease the risk of viral transmission. This article shows a new technique, which effects real-time data on COVID-19 contagion rates, indications of patients, ambulance obtainability, and geographical info. In this case, we clarify how by adding these features to the Kruskal algorithm, we reach the finest paths for the best dispatch of ambulances with a perspective of minimizing the spread of the virus. The revised algorithm is designed to focus on routes with the fewest stops and contact points, enabling faster response times and reducing opportunities for disease spread. Furthermore, we present a case study demonstrating the effectiveness of our adapted algorithm in a simulated emergency reply scenario during the COVID-19 pandemic. Comparative analysis with traditional path-shortening algorithms showcases the superior concert of our approach in terms of reply time and safety. This shows delivers a valued effect to the field of emergency medical services, offering a practical solution to address the tests posed by COVID-19 while improving overall response efficiency. The modified algorithm can be seamlessly integrated into existing ambulance dispatch systems, helping healthcare providers and emergency responders make informed decisions that prioritize both patient well-being and public health safety during these unprecedented times.

Keywords: COVID-19, shortest path, Kruskal algorithm, linked lists, minimum spanning tree

1. INTRODUCTION

The shortest route to reducing time and cost must be identified. There are multiple routes in the cocept, mentioned to as edges, which consist of the value or length of the path and are connected to the

edges of the edges that represent the cities. Its uses are found in various fields such as Google Maps, image processing, and communications. This is an example of the fact that we have a range of cities and a request to create a network at the lowest cost. In the beginning, we must select the straight route in order to attach all the cities without creating a closed circular track. This condition is similar to that of the seller's man, but it is less complex.

The shortest-spanning tree can be used for causal the optimal route. The Kruskal procedure is one of the algorithms used to control the shortest-spanning tree. It splits the tree into clusters, as well as each cluster containing two verses that are linked by a solo edge. The fundamental rule of Kruskal's algorithm is to select the edge with the minimum weight. A circular closed path is not possible due to the designated edges [1].

Emergency medical facilities faced their most extraordinary challenge through the COVID-19 pandemic. Ambulance routing systems designed for traditional purposes might prove inadequate when managing public health crises that display high infection rates through diseases. Healthcare patients must be transported fast to facilities as exposure reductions are equally vital while they move between locations.

Emergency medical facilities can use their conventional ambulance routing procedure under typical operations yet this method fails to account for pandemic-specific changes that occur in service demand. Route optimization becomes complicated because of isolation zones combined with high-risk zones while also requiring individual protective equipment (PPE). Real-time patient-ambulance communication stands essential because patient transportation needs to remain immediate [2]. The patient communicates through architectural design with the ambulance as depicted in Fig.1.

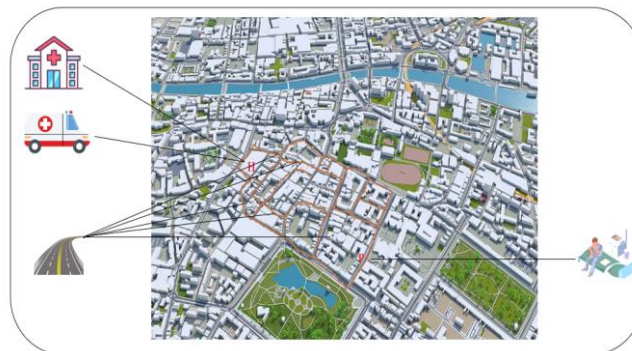


Fig. 1 Architectural communication between the patient and ambulance

1.1 The Role of the Modified Kruskal Algorithm

The Kruskal algorithm succeeds in identifying minimum spanning trees through efficient operations which when adapted can resolve communication issues in ambulance services during pandemic situations. This adapted algorithm enables the development of a safety-conscious scheme that optimizes ambulance routes operation through real-time patient need assessment[3].

A modified procedure integrates multiple criteria such as ambulance patient proximity and PPE availability along with risk assessment on different areas and emergency requirements. The system employs real-time communication between patient role and the ambulance service to create informed decisions about selecting routes.

1.2 Key Objectives of the Modified Algorithm

The modified Kruskal algorithm has several key purposes which are detailed below:

1. The Ambulance Route Optimization process establishes a path which shortens over time to lower the duration needed for ambulance transportation of a patient. The solution combines traffic data and patient conditions with road restriction information to create shortest possible routes.
2. Safety First: Prioritizing patient and ambulance personnel safety. The algorithm analyzes both infectious outbreak areas and disease risk zones and the availability of protective equipment to reduce the possibility of disease exposure.
3. Real-Time Communication functions as part of the method to keep patients and ambulance services connected in real time. The system can arrange emergency cases according to their urgency levels by tracking up-to-date patient information from patients.
4. Proper distribution of resources stands as essential for managing a pandemic situation. Through its algorithm the program directs ambulances to areas where their presence is essential first so it reduces unneeded journeys yet maximizes operational resources.

This modified Kruskal algorithm functions as a data-based method to configure ambulance routes within pandemic conditions. The system employs technological capabilities together with real-time communication to combine speed with safety measures. [4],[5], ensuring that patients receive prompt medical attention while minimizing the risk of infection transmission.

This article examines the detailed workings of the modified algorithm and implementation with focus on its impact on ambulance communication systems while COVID-19 persists. The paper is structured as following. Section II describes research regarding the Kruskal algorithm.

The shortest path is in Section III. Kruskal's algorithm in Section IV, Modified Kruskal's algorithm in Section V, and Section VI reveal the conclusion. Finally, the paper concludes with Section VII references.

2. RELATED WORK

The researchers presented a study on the use of Kruskal Algorithm Location in Work-flow Push System. An appropriate payment technology algorithm was designed. The work involved the study of the general structural characteristics of the work network on social networking sites, using the dual stack protocol, with a focus on the shortest-path location algorithm.

Drew a minimalist concept tree-spanning algorithm. It concluded with satisfactory results [2].

A study on improved delay management through mobile computing was presented by researchers. The study entailed the submission of two proposals. The first proposal suggested a new model of a three-layer network architecture. CENAM, an abbreviation for Cloud computing and edge computing, was combined.

In the measured edge layer, we propose a Calculator Mutual cooperation scheme among the edge devices and the Kruskal algorithm for calculating the minimum span weighted undirected graph tree composed of edge nodes. The results resolved the cloud computing delay issue [1].

The article discusses a study on industrial wireless sensor networks that are used by synthetic bee colony. The study is based on a tree-building algorithm. Researchers combined quantum computing with an artificial bee colony, as well as designing the construction of an extended tree for bee quantum calculations. In order to simulate the algorithm, the food source substitution strategy was improved. The proposed algorithm improved research performance and quality, according to the results [3].

The power distribution network acquisition of the minimum power flow path was explained through Kruskal's algorithm during 2019. A performance evaluation took place between Kruskal's Method and Binary Particle Swarm Optimisation (BPSO) in solving network reconfiguration problems. The current research team applies the proposed load restoration method to the 33-bus IEEE system. Observations from the obtained results demonstrate that Kruskal's algorithm runs faster than both Binary Particle Swarm Optimisation (BPSO) [4].

3. SHORTEST PATH

When searching the shortest path in an undirected graph (G) from vertex i to destination vertex j one needs to select the sequences of vertices that will become the second then subsequent vertices until reaching j . A path can be minimized through the selection of the smallest side. The following Fig. 2 displays an undirected graph referred to as a digraph. An examination of the minimum pathway exists between vertex 1 and vertex 3

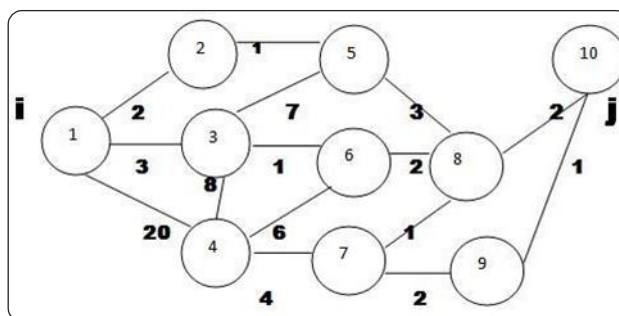


Fig. 2. Undirected graph G

4. KRUSKAL ALGORITHM

Kruskal's algorithm is one of the algorithms that finds the shortest spanning tree for a connected weighted graph [2]. The Minimum Spanning Tree (MST) is a sub-graph of the graph $G=(V, U)$. It includes all the vertices with the least cost of undirected and weighted edges and has the minimum total edge weight [3-7]. Kruskal's algorithm is based on segmenting the graph into clusters and arranging them in ascending order. The weight of the edges determines the priority in the queue [8-9]. It can also be applied in the fields of education for choosing the shortest path for students and in communications for reducing costs in establishing conveyor roads [10-11]. A reviewed the number of studies on the Kruskal algorithm [12-15], Path-Planning Algorithms [16], and Prim's and Boruvka's algorithm [17]. Recently, studies have discussed the issue of communication between patients and ambulances in relation to the speed of receiving treatment, particularly during the COVID-19 pandemic. These studies include references [18-31].

4.1 WORK OF KRUSKAL'S ALGORITHM

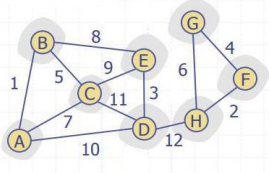
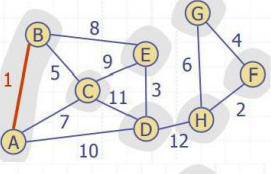
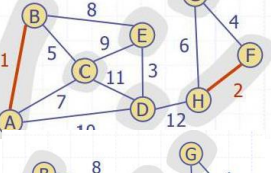
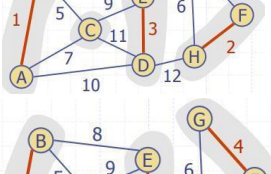
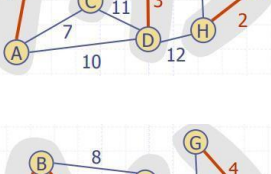
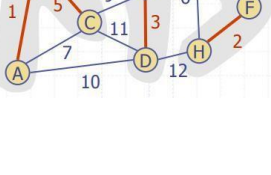
- 1) Divide the graph into clusters and sort them in ascending order.
- 2) choose the shortest edge in a graph.
- 3) Next, select the shortest edge.

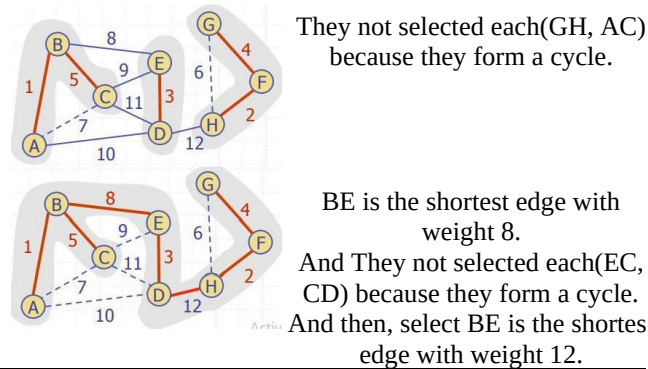
Which one does not make a cycle? Repeat step 2 until all vertices have been connected. Choose an edge that has the lowest cost and does not create a cycle when included in the set of edges that have already been chosen. Note: Input the edge weight into the algorithm, and the output is the shortest spanning tree [2]. The pseudo code is shown below, and Table 1.

```

MST-Kruskal ( $G, w$ )
1  $A \leftarrow \emptyset$ 
2 for each vertex  $v \in V[G]$  do
3   Make-Set ( $v$ )
4 sort the edges of  $E$  by nondecreasing weight  $w$ 
5 for each  $(u, v) \in E$ , in nondecreasing of weight do
6   if Find-Set ( $u$ )  $\neq$  Find-Set ( $v$ ) then
7      $A \leftarrow A \cup \{(u, v)\}$ 
8     Union (Set ( $u$ ), Set ( $v$ ))
9 return  $A$ 
  
```

Table 1. Work of kruskal's algorithm

Image	Edge case description
	Example to Kruskal's algorithm
	AB is the shortest edge with weight 1.
	HF is the shortest edge with weight 2.
	DE is the shortest edge with weight 3.
	GF is the shortest edge with weight 4.
	BC is the shortest edge with weight 5.



4.2 THE RUNTIME FOR THE ALGORITHM KRUSKAL'S

Kruskal's algorithm consists of two stage:

The first step is to initialise A that takes time $O(1)$.

The second step is to sort the clusters and takes time $E \log E$ [5]. Performing: loop takes time $|E|-1$.

$|V|$: MakeSet ().

$|E|$: FindSet ().

$V|-1$: Union ().

The total running time is $E \log E$.

5. MODIFICATION OF THE KRUSKAL'S ALGORITHM

The idea of modification was to reduce the period complexity of this algorithm, which depends on its sorted state. It is known that the time taken for sorting is $E \log E$. Currently, there was no research algorithm that is less than this time in paper. We reduced the $E-1$ iteration time by finding the last edge without reaching an end. The technique we used was Linked List.

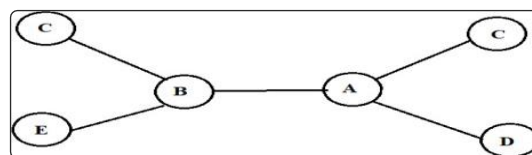


Fig. 3. Nodes network

We modified the MakeSet() function to return the amount of vertices and the grade of each vertex. And the contents of each vertex and a discover variable that specifies whether or not the node was detected. Fig. 3 and Fig. 4 displays the variables the node.



Fig. 4. The components of the node

5.1 Pseudocode of Modification Kruskal's Algorithm

Algorithm $\text{KruskalMST}(G)$:

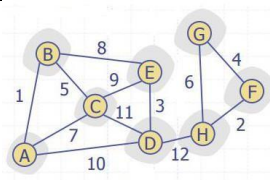
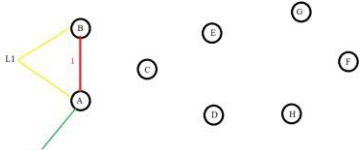
Input: A simple connected weighted graph G with n vertices and m edges
Output: A minimum spanning tree T for G

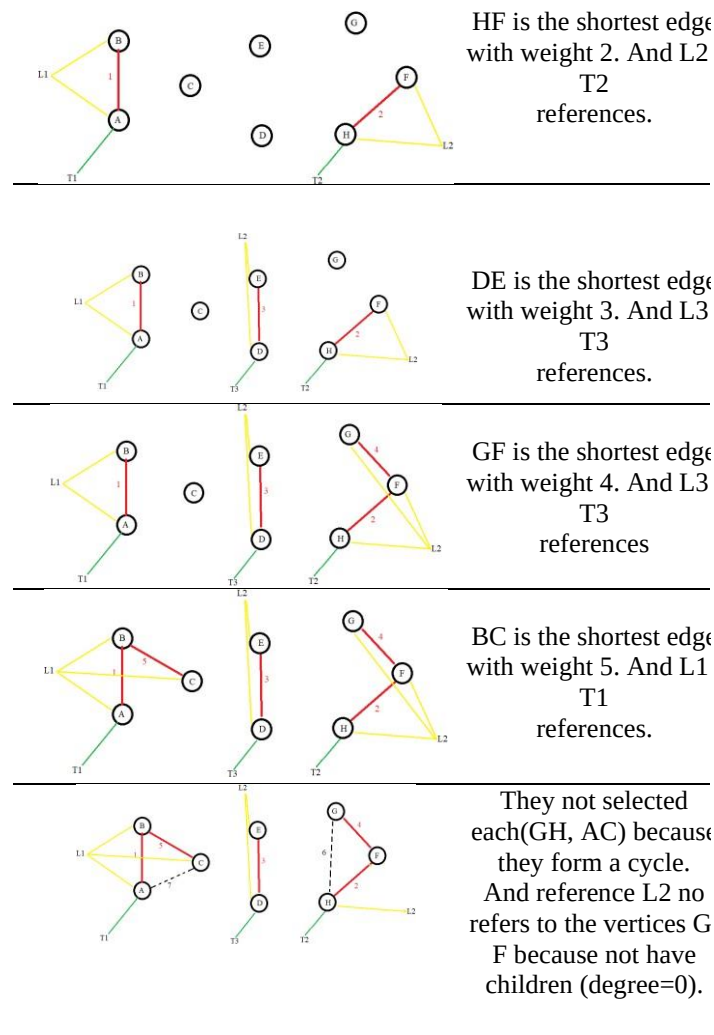
for each vertex v in G **do**
 Define an elementary cluster $C(v) \leftarrow \{v\}$.
Let Q be a priority queue storing the edges in G , using edge weights as keys
 $T \leftarrow \emptyset$ // T will ultimately contain the edges of the MST
while T has fewer than $n - 1$ edges **do**
 $(u, v) \leftarrow Q.\text{removeMin}()$
 Let $C(v)$ be the cluster containing v
 Let $C(u)$ be the cluster containing u
 if $C(v) \neq C(u)$ **then**
 Add edge (v, u) to T
 if v or u not Discovered **then**
 $v = \text{discover}$ or $u = \text{discover}$
 Link = children v and u
 vertex = vertex + 1
 Merge $C(v)$ and $C(u)$ into one cluster, that is, union $C(v)$ and $C(u)$
 if $T=2$ and all vertices is discovered **then**
 break
 $T = \min(\text{degree}(\text{Link1}), \text{degree}(\text{Link2}))$
 edge = SearchEdge(T) // search connect edge between $T1$ & $T2$ from T
 $T = \text{Merge}(T1, T2)$ will merge $T1$ and $T2$ with last edge
Return T

5.2 Work of our_Kruskal's Algorithm

Initially, the modified algorithm behaves exactly like the classic algorithm. Each cluster has two references until we get two clusters; the first reference refers to the top of the cluster to know the circle when adding. The other reference is to the end of the vertices that contain children, Kruskal's algorithm stopped functioning at that step because there were two trees that contained all the nodes in Fig. 5. The stop of Kruskal's is shown below.

Table 2. Work of our_kruskal's algorithm

Image	Edge case description
	example to Kruskal's algorithm.
	AB is the shortest edge with weight 1. And L1 T1 references.



The references T1 and T2 pertain to all vertices within the same tree. L1 and L2 refer to all vertices that contain edges and are not detected as belonging to the same tree (i.e., they have degrees that are not equal to zero).

The reference L1 was compared with L2, whichever contained the lowest degree was searched. In Table 2, L2 had the smallest degree. On L2, the last edge connecting the trees T1 and T2 was searched.

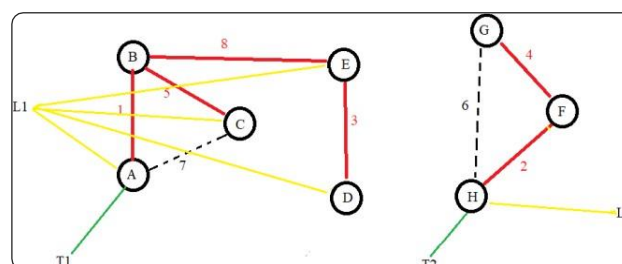


Fig. 5. Steps of Kruskal's

5.2 The runtime for the algorithm Our_Kruskal's

Our Kruskal's algorithm consists three stage.

The first step is to initialise A that takes time $O(1)$.

The second step is to sort the clusters and takes time $E \log E$.

performing:

$|V|$: MakeSet() .

$|E|$: FindSet().

$|V|-1$: Union().

We have two trees L_1, L_2

$L = \text{minimum}(\text{deg}(L_1), \text{deg}(L_2))$ take time $O(1)$. Search(L) take time $O(\text{deg}(L))$.

Loop takes time $|E| - \text{deg}(L_1)$, such that $\text{deg}(L_1)$ is from 1 to $E / 2$.

The total running time is $E \log E$.

The interpretation of the results (Basic concept of the classical and modified algorithms and Comparison of the classical and modified algorithms) is shown in Table 3 and Table 4 respectively.

Table 3. Basic concept of the classical and modified algorithms

No.	Classical Kruskal Algorithm	Modified Kruskal Algorithm
1	The classical Kruskal algorithm is used for finding the minimum spanning tree in a connected weighted graph.	The modified Kruskal algorithm aims to reduce the time complexity of the classical algorithm by avoiding unnecessary edges.
2	It operates by selecting the edges with the smallest weights while ensuring that no closed circular paths are formed.	It uses linked lists to manipulate the edges and neglect those that do not contribute to the minimum spanning tree.
3	The algorithm's time complexity is $O(E \log E)$, where E is the number of edges.	The time complexity of the modified algorithm is still $O(E \log E)$ but may be more efficient in practice due to edge elimination

Table 4. Comparison of the classical and modified algorithms

No.	Function	Classical Kruskal Algorithm	Modified Kruskal Algorithm
1	Efficiency	The time complexity of $O(E \log E)$	Reducing the number of edge calculations
2	Edge Elimination	Edge is present	Edge is not present
3	Application	It finds Minimum Spanning Trees, benefiting network design, clustering, circuit design, transportation, biology, image processing, and more.	It considers factors like proximity to patients, availability of personal protective equipment (PPE), quarantine zones. real-time communication.
4	Flexibility	Can not change	It is adaptable to changing conditions. evolving data
5	Safety	high	high

6. CONCLUSION

In conclusion, the modified Kruskal algorithm presented in this study represents a significant advancement in optimizing communication pathways between patients and ambulances during the COVID-19 pandemic. The algorithm's integration of real-time COVID-19 data, patient symptoms, ambulance availability, and geographical information offers a comprehensive and adaptable solution for emergency response systems. This study demonstrates the practical benefits of this algorithm through a compelling case study, which showcases its superior performance in terms of reducing response times and minimizing the risk of viral transmission. These findings underscore the algorithm's potential to revolutionize the way emergency services operate during public health crises. By prioritizing paths that minimize stops and contact points, our algorithm ensures a faster and safer response to patients in need, which is crucial during a pandemic where every moment counts. Furthermore, its seamless integration into existing ambulance dispatch systems allows for efficient adoption and implementation. Thus, the modified Kruskal algorithm for shortening communication pathways between patients and ambulances during COVID-19 represents a promising innovation in the field of emergency medical services. It empowers healthcare providers and emergency responders to make data-driven decisions that prioritize patient well-being and public health safety, ultimately saving lives and reducing the spread of infectious diseases in the face of unprecedented challenges.

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